

PHY682 Special Topics in Solid-State Physics: Quantum Information Science

Lecture time: 2:40-4:00PM Monday & Wednesday

Today 10/14:

1. More on topological quantum computation:
Magic state distillation and surface code quantum computation

Week 8: More topological
please: Topological quantum
computation continued,
magic state distillation, and
surface code

Arbitrary phase gate from 'phase' state

With the state available

$$|A_\theta\rangle = (|0\rangle + e^{i\theta}|1\rangle)/\sqrt{2}$$

we can implement a general phase gate $U_\theta = \begin{pmatrix} 1 & 0 \\ 0 & e^{i\theta} \end{pmatrix}$

Suppose we have $|\psi\rangle \otimes |A_\theta\rangle = (a|0\rangle + b|1\rangle) \otimes (|0\rangle + e^{i\theta}|1\rangle)/\sqrt{2}$

*want to apply U_θ ?
assume Clifford gates are
of high fidelity*

we want to measure stabilizer $Z \otimes Z$ (+1 eigenspace spanned by 00, 11; -1 eigenspace by 01, 10)

$$|\psi\rangle \otimes |A_\theta\rangle = \frac{1}{\sqrt{2}} [(a|00\rangle + e^{i\theta}b|11\rangle) + e^{i\theta}(a|01\rangle + e^{-i\theta}b|10\rangle)]$$

Each with probability 1/2

+1 eigenspace *-1 eigenspace*

+1 (i) $a|00\rangle + e^{i\theta}b|11\rangle$

CNOT

$$(a|0\rangle + e^{i\theta}b|1\rangle) \otimes |0\rangle = U_\theta |\psi\rangle \otimes |0\rangle$$

*apply $U_{-\theta}$
gate
random (+1) (-1)*

-1 (ii) $a|01\rangle + e^{-i\theta}b|10\rangle$

$$(a|0\rangle + e^{-i\theta}b|1\rangle) \otimes |1\rangle = U_{-\theta} |\psi\rangle \otimes |1\rangle$$

Magic states $|H\rangle$ and $|T\rangle$

Clifford gate G_c
 G_c Pauli: $G_c^{-1} = \text{Pauli}$

□ If we have $|H\rangle = \cos\left(\frac{\pi}{8}\right)|0\rangle + \sin\left(\frac{\pi}{8}\right)|1\rangle$ $U_\theta = \begin{pmatrix} 1 & 0 \\ 0 & e^{i\theta} \end{pmatrix}$ $S = \begin{pmatrix} 1 & 0 \\ 0 & e^{i\pi/2} \end{pmatrix}$

Clifford

then $(HS)|H\rangle = (e^{i\pi/8}|0\rangle + e^{-i\pi/8}|1\rangle)/\sqrt{2} = e^{i\pi/8}|A_{-\pi/4}\rangle \Rightarrow$ non-Clifford

→ Can implement non-Clifford gate $T^\dagger = U_{-\pi/4}$

$$T = \begin{pmatrix} 1 & 0 \\ 0 & e^{i\pi/4} \end{pmatrix}$$

use T to denote the "T" gate in the paper by Bravyi & Kitaev

□ If we have $|T\rangle = \cos\beta|0\rangle + e^{i\pi/4}\sin\beta|1\rangle$, $\cos(2\beta) = 1/\sqrt{3}$

can we also generate a useful phase state? \Rightarrow from that I can get a phase gate

➤ Start with $|\Psi_0\rangle = |T\rangle \otimes |T\rangle$ and measure stabilizer $Z \otimes Z$

Only keep +1 outcome (probability 2/3)

When it succeeds \Rightarrow

$$\cos^2\beta|00\rangle + e^{i\pi/2}\sin^2\beta|11\rangle \sim |\Psi_1\rangle = \cos\left(\frac{\pi}{12}\right)|00\rangle + i\sin\left(\frac{\pi}{12}\right)|11\rangle$$

$$(\cos\beta|0\rangle + e^{i\pi/4}\sin\beta|1\rangle) (\cos\beta|0\rangle + e^{i\pi/4}\sin\beta|1\rangle) U_{-\pi/6}$$

non-Clifford

➤ Apply CNOT and then Hadamard: first qubit becomes $(e^{i\pi/12}|0\rangle + e^{-i\pi/12}|1\rangle)/\sqrt{2} = e^{i\pi/12}|A_{-\pi/6}\rangle$

[Bravyi & Kitaev, PRA71, 022316 (2005)]

Magic state distillation

applied to Ising anyons

How do we generate high-fidelity H and T states?

$$|H\rangle = \cos\left(\frac{\pi}{8}\right)|0\rangle + \sin\left(\frac{\pi}{8}\right)|1\rangle \quad H|H\rangle = |H\rangle$$

$$|T\rangle = \cos\beta|0\rangle + e^{i\pi/4}\sin\beta|1\rangle, \quad \cos(2\beta) = 1/\sqrt{3}$$

$$\tilde{T} = e^{i\pi/4}SH \quad \tilde{T}|T\rangle = e^{i\pi/3}|T\rangle \quad |T^\perp\rangle \equiv \sigma_y H|T\rangle \quad \tilde{T}|T^\perp\rangle = e^{-i\pi/3}|T^\perp\rangle$$

Bravyi & Kitaev exploited the 5-qubit error correcting code for T state and CSS codes for H state.

[complicated procedure]

→ An iteratively procedure that takes many noisy copies to distill a single high-quality T and H states, provided, respectively

$$\epsilon_T \equiv 1 - \langle T | \rho_{\text{noisy}} | T \rangle < \epsilon_0 \approx 0.173$$

$$\epsilon_H \equiv 1 - \langle H | \rho_{\text{noisy}} | H \rangle < \epsilon_0 \approx 0.141$$

Summary:
use noisy gates to create states close to $|T\rangle$ & $|S\rangle$
⇒ can distill clean $|T\rangle$ & $|S\rangle$

close



↓
can implement high-fidelity non-Clifford gates

+ "protected" Clifford gate
⇒ protected QC

[Bravyi & Kitaev, PRA71, 022316 (2005)]

T and H type states on Bloch sphere

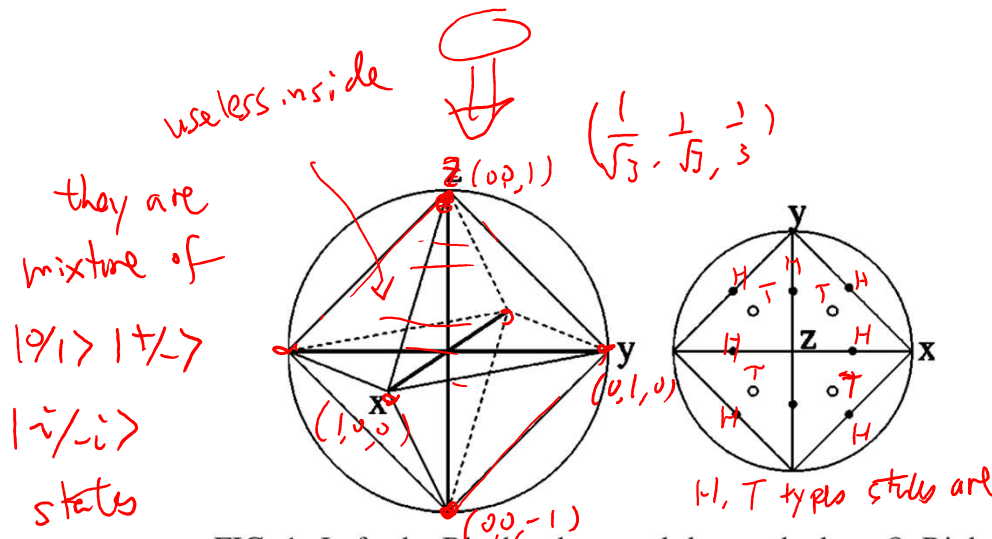


FIG. 1. Left: the Bloch sphere and the octahedron O . Right: the octahedron O projected on the x - y plane. The magic states correspond to the intersections of the symmetry axes of O with the Bloch sphere. The empty and filled circles represent T -type and H -type magic states, respectively.

$$|T\rangle\langle T| = \frac{1}{2} \left[I + \frac{1}{\sqrt{3}} (\sigma^x + \sigma^y + \sigma^z) \right]$$

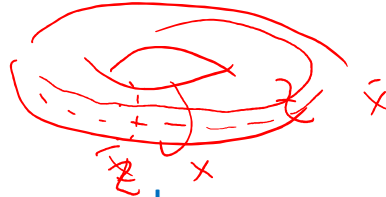
$$|H\rangle\langle H| = \frac{1}{2} \left[I + \frac{1}{\sqrt{2}} (\sigma^x + \sigma^z) \right]$$

$$\left(\frac{1}{\sqrt{2}}, 0, \frac{1}{\sqrt{2}} \right)$$

Clifford gates simply ~~move~~ rotate $|T\rangle\langle T|$ & $|H\rangle\langle H|$

Surface code

Bravyi & Kitaev, quant-ph/9811052
 Raussendorf, Harrington & Goyal, NJP 9, 199 (2007)
 Folwer et al., Phys. Rev. A 86, 032324 (2012)
 Fujii, arXiv:1504.01444



It is essentially the toric code on a planar geometry

Star operators:

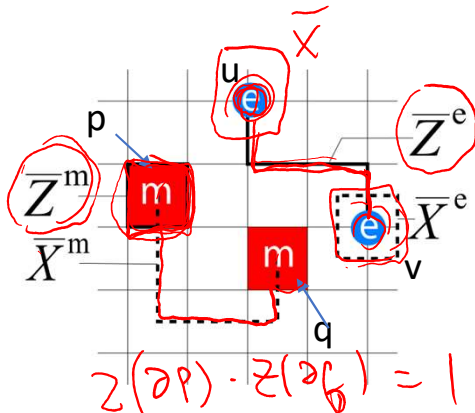
$$A_s = \prod_{j \in s} \sigma_x^{[j]} \quad \begin{array}{c|c} \sigma_x & \sigma_x \\ \hline \sigma_x & \sigma_x \end{array}$$

Plaquette operators:

$$B_p = \prod_{j \in \partial(p)} \sigma_z^{[j]} \quad \begin{array}{c} \sigma_z \\ \sigma_z \quad \square \quad \sigma_z \\ \sigma_z \end{array}$$

remove the two stabilizers
 $-2 + 1 = -1$
 \Rightarrow one qubit
 remove constraints

But a logical qubit is encoded in a pair of defects (or holes)



Two kinds:

(1) a pair of plaquette defects $\langle p, q \rangle \rightarrow B_p$ and B_q no constraint but $B_p B_q$ is in stabilizer

(2) a pair of star (vertex) defects $\langle u, v \rangle \rightarrow A_u$ and A_v no constraint but $A_u A_v$ is in stabilizer

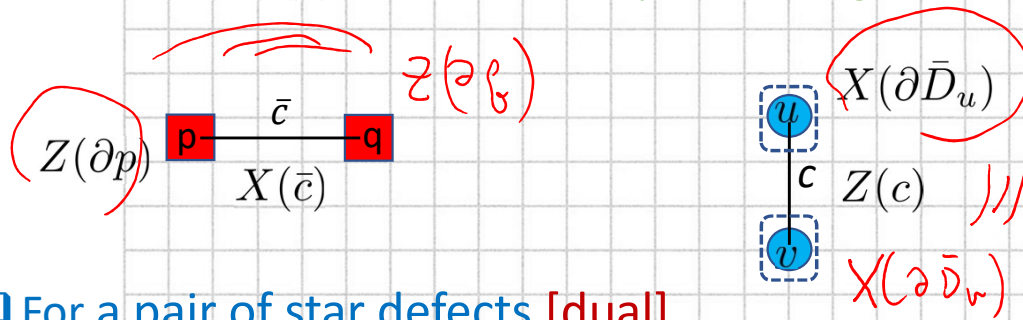
remove the two stabilizers
 $-2 + 1 = -1$
 \Rightarrow one qubit
 remove constraints

Logical operators

□ For a pair of plaquette defects [primal]

Logical Z: $Z(\partial p) \equiv Z(\partial q) \rightarrow$ Product of Z operators around plaquette p or q

Logical X: $X(\bar{c}) \rightarrow$ Product of X operators along line \bar{c}



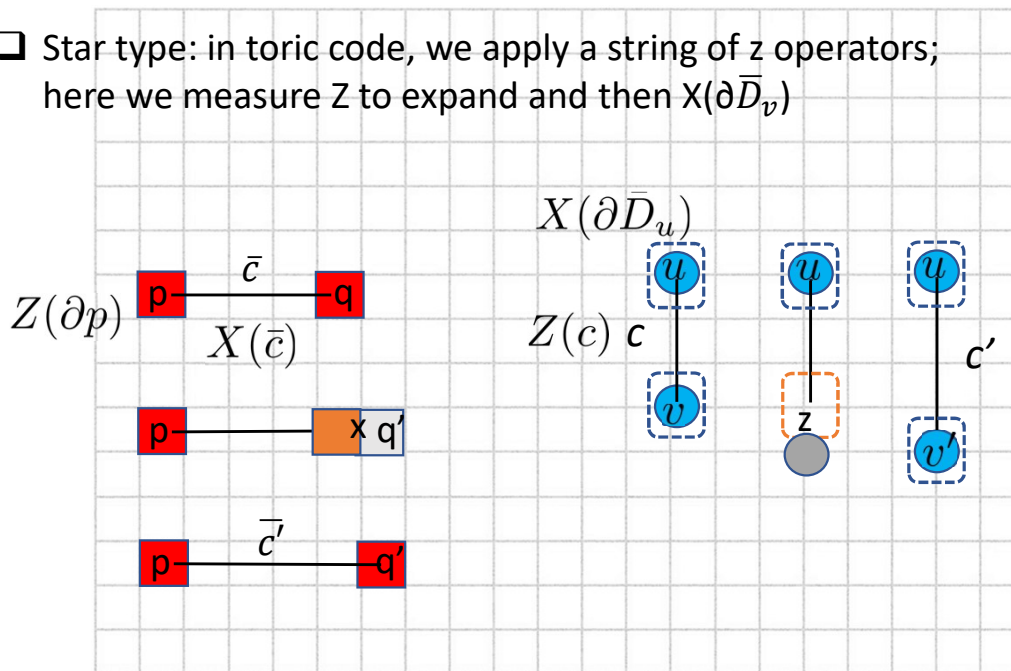
□ For a pair of star defects [dual]

Logical Z: $Z(c) \rightarrow$ Product of Z operators along line c

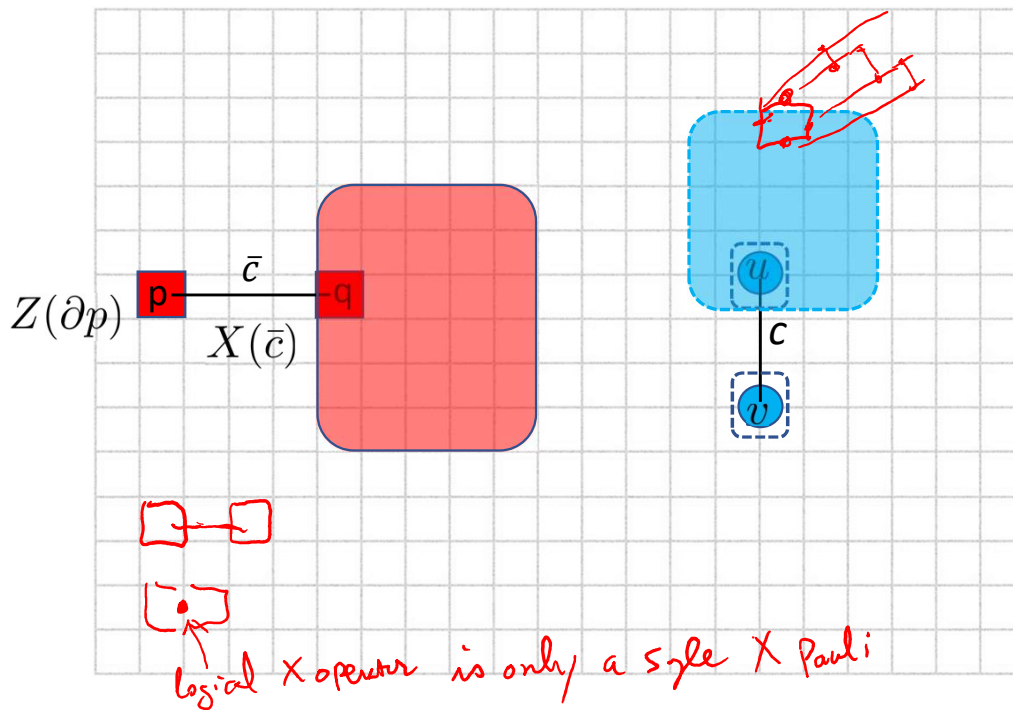
Logical X: $X(\partial \bar{D}_u) \equiv X(\partial \bar{D}_v) \rightarrow$ Product of X operators around vertex u or vertex v

Logical operators: deformable

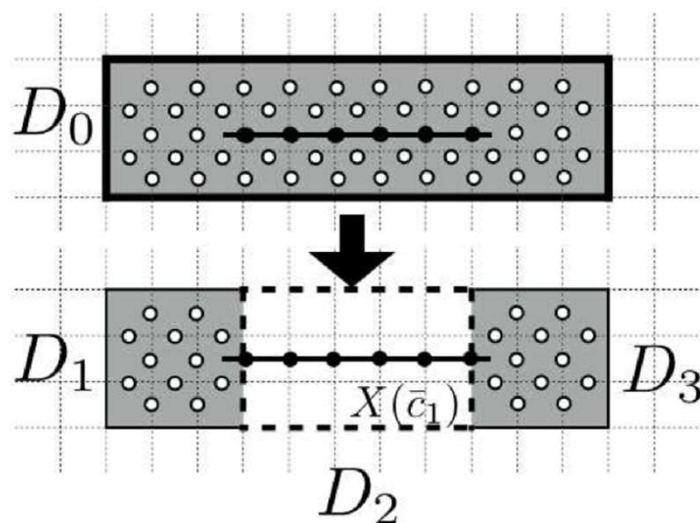
- ❑ Plaquette type: in toric code, we apply a string of x operators; here we measure X to expand and then $Z(\partial q)$
- ❑ Star type: in toric code, we apply a string of z operators; here we measure Z to expand and then $X(\partial \bar{D}_v)$



Logical operators: expandable/shrinkable

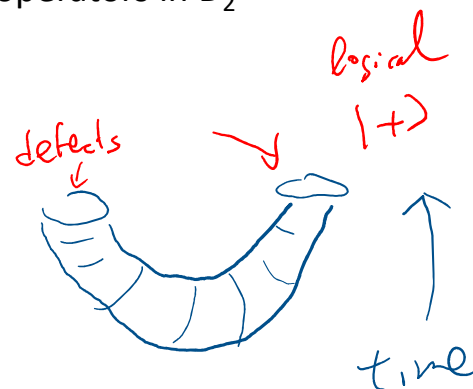


Logical X-basis state preparation



prepare logical $|+\rangle$ → break stabilizer operators in that region

1. Create a defect region D_0 by measuring all qubits (circles) in X basis $|+\rangle \text{ or } |-\rangle \Rightarrow |+\rangle$
2. Annihilate defect region D_2 by measuring star operators in D_2

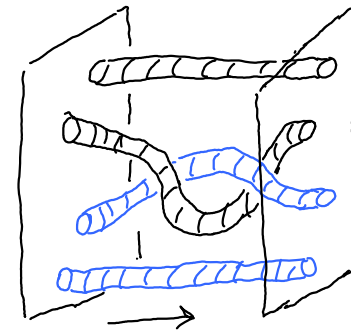
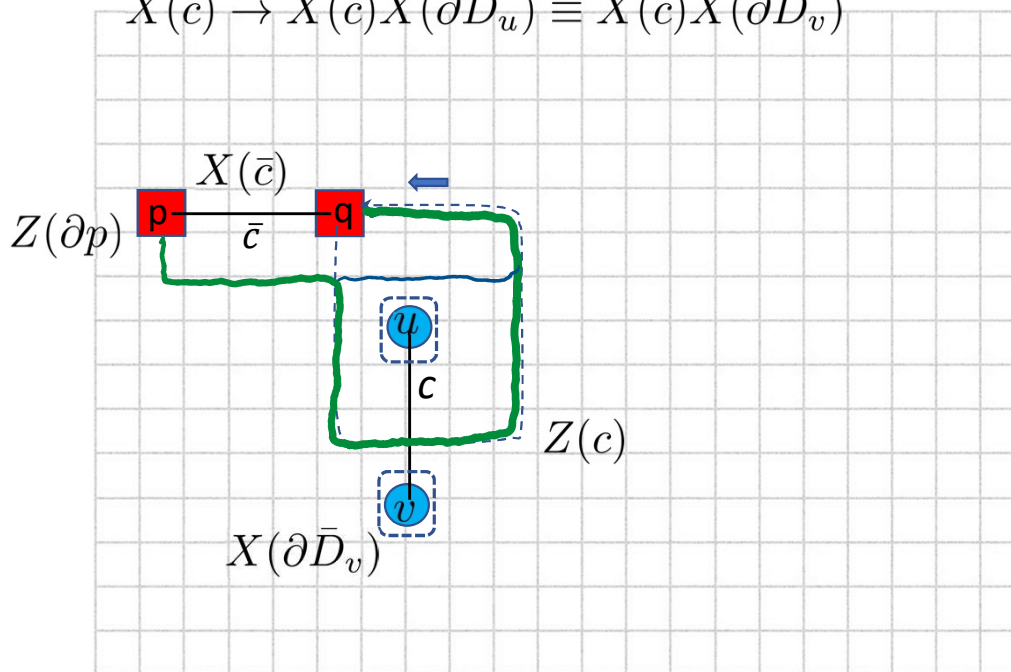


Protected CNOT gate by braiding: m around e

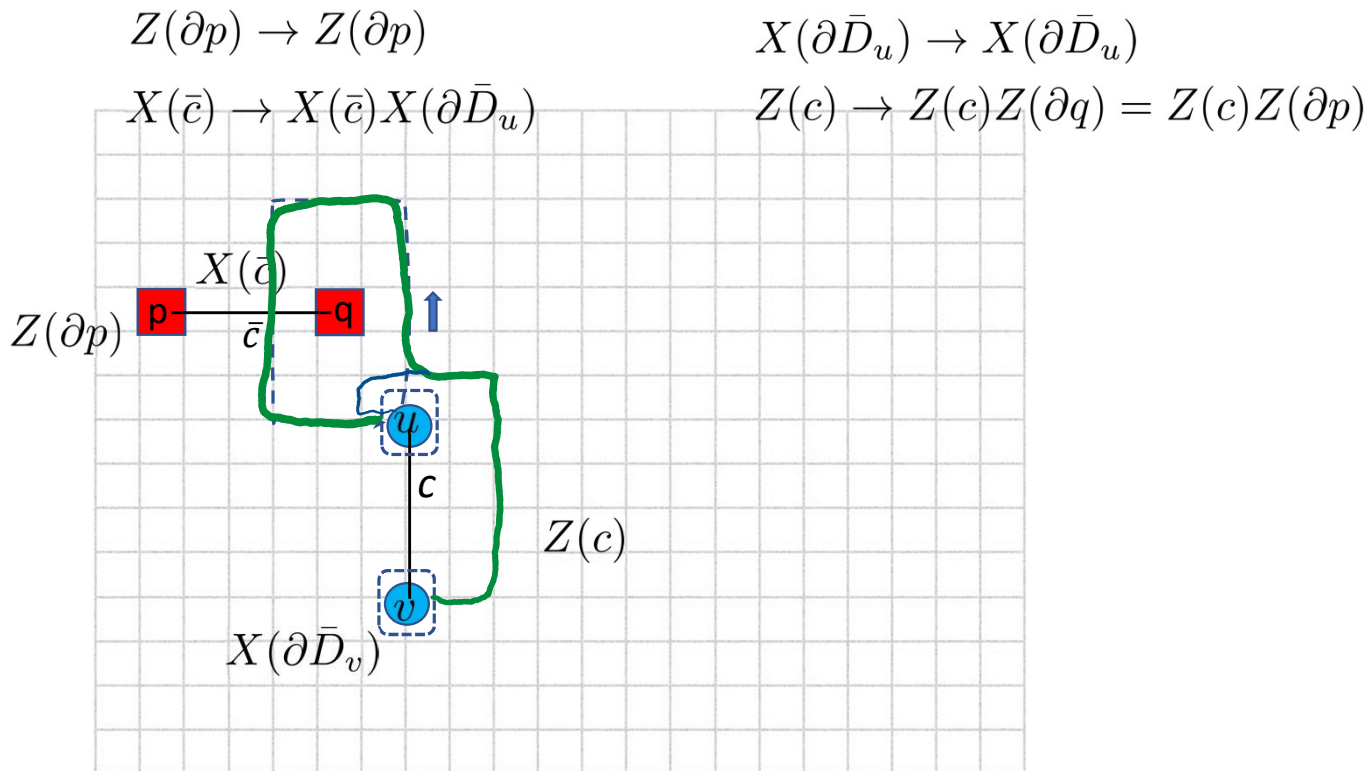
$$Z(\partial p) \rightarrow Z(\partial p)$$

$$X(\partial \bar{D}_u) \rightarrow X(\partial \bar{D}_u)$$

$$X(\bar{c}) \rightarrow X(\bar{c})X(\partial \bar{D}_u) \equiv X(\bar{c})X(\partial \bar{D}_v)$$

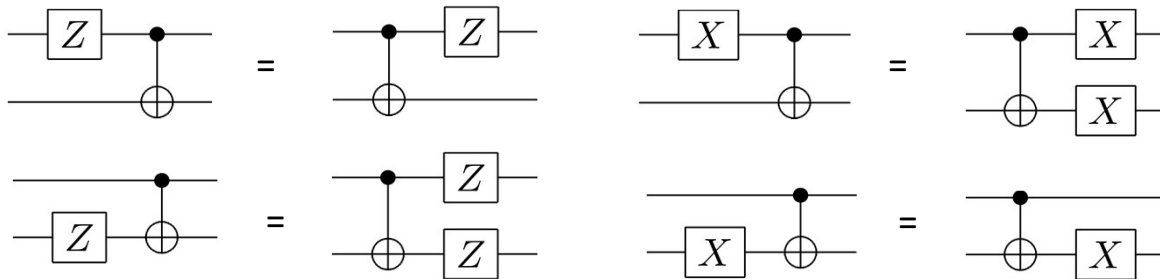


Equivalent view: braid e around m



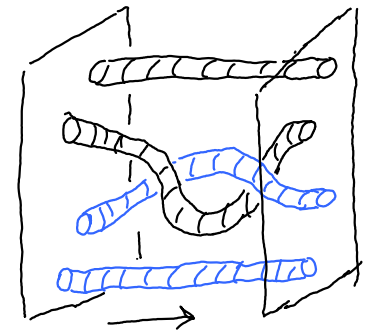
Verifying CNOT

□ Recall:



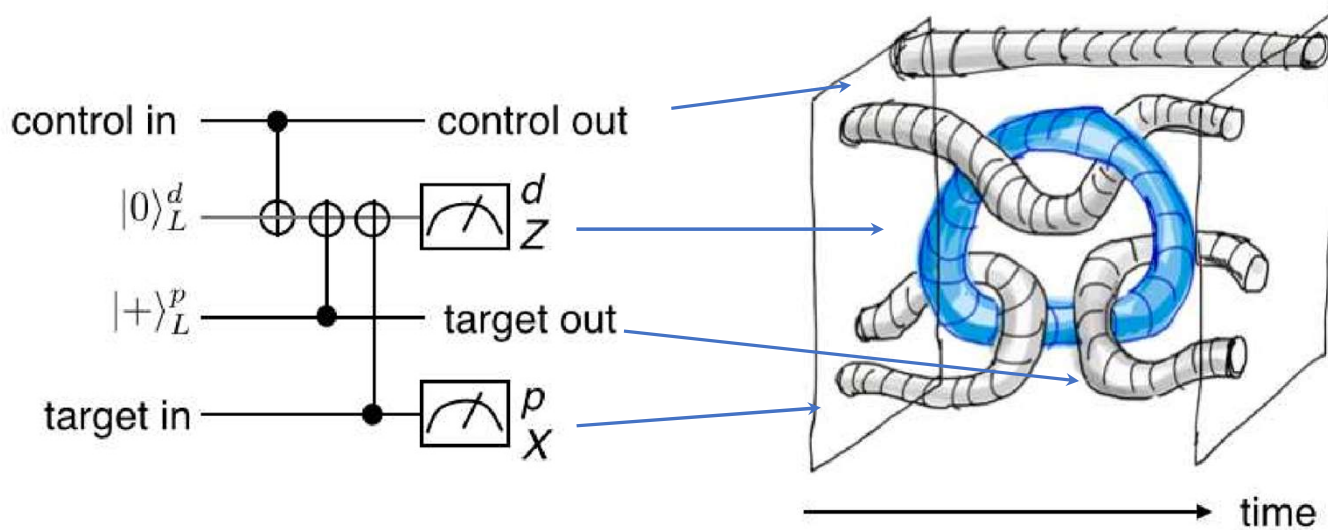
□ We have:

$$\begin{aligned}
 Z(\partial p) &\rightarrow Z(\partial p) & X(\partial \bar{D}_u) &\rightarrow X(\partial \bar{D}_u) \\
 X(\bar{c}) &\rightarrow X(\bar{c})X(\partial \bar{D}_u) & Z(c) &\rightarrow Z(c)Z(\partial q) = Z(c)Z(\partial p)
 \end{aligned}$$



→ Primal qubit is the *control* and dual qubit is the *target*

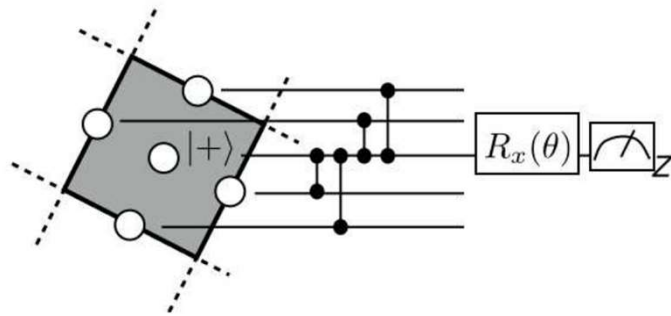
CNOT gate between two “primal” pairs



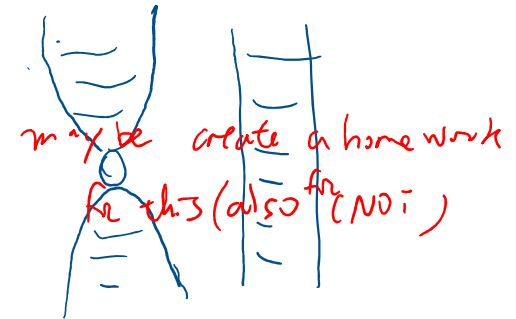
Fujii: arXiv 15...

Logical Z and X rotations (unprotected)

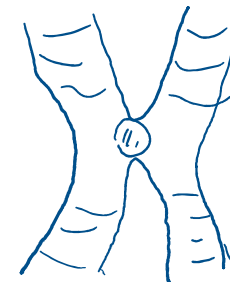
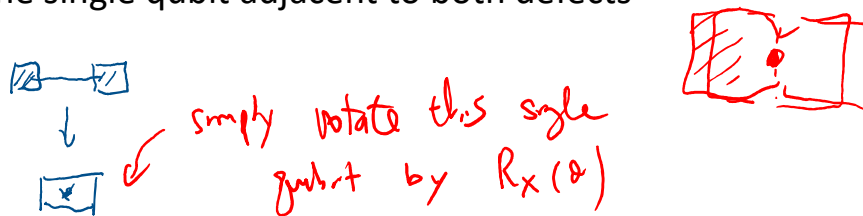
- **Logical Z rotation:** by adding an ancilla in + state, perform CZ gates (below), rotate the ancilla before Z measurement



$$e^{-i(-1)^m(\theta/2)Z(f_i)}$$



- **Logical X rotation:** bring two defects next to each other and rotate directly on the single qubit adjacent to both defects

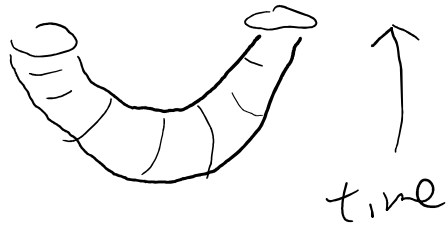


Even though they are not protected, they can be used to inject noisy magic states → then distill to a higher-fidelity magic state

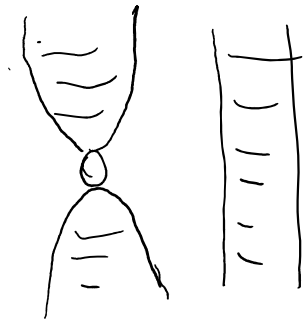
Logical arbitrary phase state

$$\frac{1}{\sqrt{2}}(e^{-i\theta/2}|0\rangle + e^{i\theta/2}|1\rangle) = e^{-i\theta Z/2}|+\rangle$$

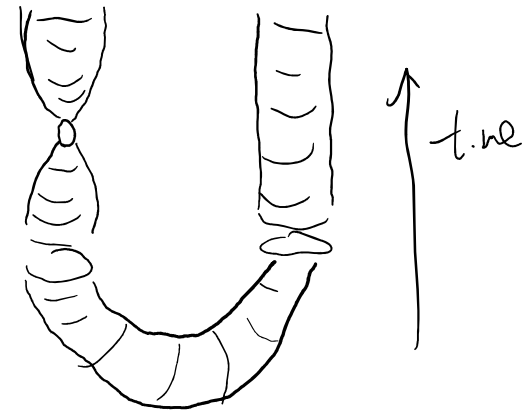
□ First prepare logical + state



□ Then perform logical Z rotation



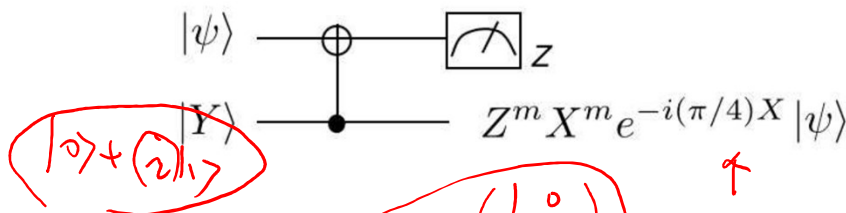
Combined picture



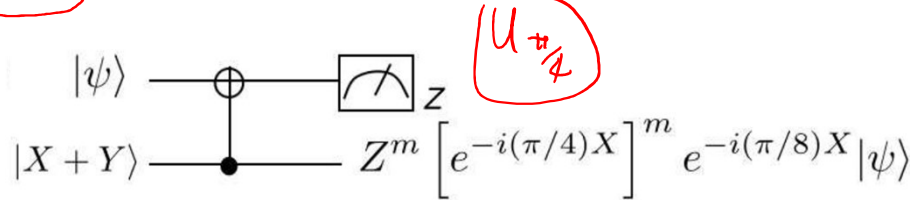
S gate and T gate

□ So far, we have a protected CNOT, but no protected single-qubit gates

- Need the eigenstate of the Pauli-Y operator, which is used to implement the **S gate** via gate teleportation



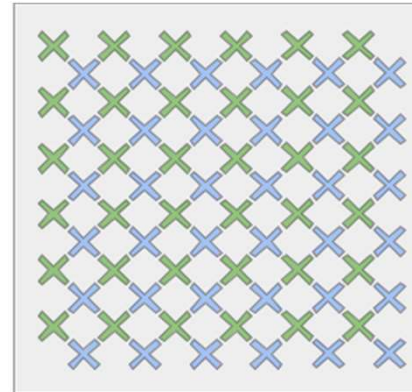
- Need the eigenstate of the $(X + Y) / \sqrt{2}$ operator, which is used to implement the **$\pi/8$ gate**, a non-Clifford gate necessary for universal quantum computation.



Surface code quantum computation

- Naturally protected CNOTs, but other gates from magic state distillation can be made of low error rates → Universal
- Google is heavily invested on surface code and built devices suitable for its implementation

(Can also apply magic state distillation to Ising anyons to make universal gates)
Microsoft



[from Google]